Semi-Federated Scheduling of Mixed-Criticality System for Sporadic DAG Tasks

Tao Yang¹, **Yue Tang²**, Xu Jiang², Qingxu Deng¹, Nan Guan² ¹Northeastern University, China, ²The Hong Kong Polytechnic University, Hong Kong

Outline

• Background & Contributions

• Preliminaries

• Semi-Federated Mixed-Criticality Algorithm

• Evaluation

Outline

• Background & Contributions

• Preliminaries

• Semi-Federated Mixed-Criticality Algorithm

Evaluation

Mixed-Criticality Systems

Sub-systems with different criticality levels (e.g. avionics):

- Monitoring and control sub-system
 - Anti-collision sub-system
 - Navigation sub-system
 - …
- low

high

Infotainment sub-system





Characteristics :

- More complex functions of system
- No reduction in system security requirements

Challenges:

- Multicore: no enough analytical techniques
- Task parallelization: dependencies and competitions among tasks

How to assign computation resources to parallel tasks on a multicore platform while guaranteeing security ?

Propose a semi-federated mixed-criticality scheduling algorithm

• Propose a mapping algorithm: from mixed-criticality tasks to a

middleware layer (mixed-criticality container tasks)

• Prove the schedulablity of our proposed algorithm

Outline

Background & Contributions

• Preliminaries

Semi-Federated Mixed-Criticality Algorithm

Evaluation

Mixed-Criticality Task Model

- Implicit-deadline sporadic parallel
- Each task is denoted by a directed acyclic graph
 - ✓ Vertices: sequential subtasks
 - ✓ Edges: dependencies
- Two types: low-criticality and high-criticality



Mixed-Criticality Task Model

- WCET of a vertex (subtask): c^N and c^O
- WCET of a task: C^N and C^O
- The longest chain of a task: L^N and L^O
- **Deadline** of a task: *D*
- Virtual deadline of a task: D'

• **Task** utilization:
$$u^N = \frac{C^N}{D'}$$
 $u^O = \frac{C^O}{D - D'}$



Existing work: Federated Mixed-Criticality Scheduling

- Task types: LH, HVH, HMH
- System states: normal state and critical state
- Assigning strategy: independently usable cores to each task in the normal and critical states.





Outline

Background & Contributions

Preliminaries

• Semi-Federated Mixed-Criticality Algorithm

Evaluation

Architecture of Our Algorithm

Semi-Federated Mixed-Criticality Algorithm:

• Two-level hierarchical scheduling framework



Architecture of Our Algorithm

Semi-Federated Mixed-Criticality Algorithm:

- Two-level hierarchical scheduling framework
- **Top-level** calculates the mapping from mixed-criticality tasks to mixed-criticality container-tasks.



Architecture of Our Algorithm

Semi-Federated Mixed-Criticality Algorithm:

- Two-level hierarchical scheduling framework
- Top-level calculates the mapping from mixed-criticality tasks to mixed-criticality container-tasks.
- **Bottom-level** schedules mixedcriticality container-tasks on physical processors.





- Processor resources interface
- Calculate number of mixedcriticality container tasks in both normal states and critical states



- Processor resources interface
- Semi-federated: Mixed-criticality container tasks provide virtual resources, so the number of mixed-criticality task can be decimal, avoiding resource waste



- Processor resources interface
- If the WCET of a task equals **c** when it executes on a unit speed processor, then its WCET on a container-task with speed σ is: $t = \frac{c}{r}$



A mixed-criticality task



$$c_{\nu 0}^{N} = 1 c_{\nu 0}^{O} = 2, \qquad c_{\nu 1}^{N} = 1 c_{\nu 1}^{O} = 3$$
$$c_{\nu 2}^{N} = 1 c_{\nu 2}^{O} = 3, \qquad c_{\nu 3}^{N} = 1 c_{\nu 3}^{O} = 3$$
$$c_{\nu 4}^{N} = 1 c_{\nu 4}^{O} = 3, \qquad c_{\nu 5}^{N} = 1 c_{\nu 5}^{O} = 2$$
$$D = 13, D' = 5$$

A mixed-criticality task



$$c_{v0}^{N} = 1 c_{v0}^{O} = 2, \qquad c_{v1}^{N} = 1 c_{v1}^{O} = 3$$

$$c_{v2}^{N} = 1 c_{v2}^{O} = 3, \qquad c_{v3}^{N} = 1 c_{v3}^{O} = 3$$

$$c_{v4}^{N} = 1 c_{v4}^{O} = 3, \qquad c_{v5}^{N} = 1 c_{v5}^{O} = 2$$

$$D = 13, D' = 5$$

$$C^{N} = 6, C^{O} = 16, L^{N} = 3, L^{O} = 7$$

A mixed-criticality task



$$c_{\nu 0}^{N} = 1 \ c_{\nu 0}^{O} = 2, \qquad c_{\nu 1}^{N} = 1 \ c_{\nu 1}^{O} = 3$$

$$c_{\nu 2}^{N} = 1 \ c_{\nu 2}^{O} = 3, \qquad c_{\nu 3}^{N} = 1 \ c_{\nu 3}^{O} = 3$$

$$c_{\nu 4}^{N} = 1 \ c_{\nu 4}^{O} = 3, \qquad c_{\nu 5}^{N} = 1 \ c_{\nu 5}^{O} = 2$$

$$D = 13, \ D' = 5$$

$$C^{N} = 6, \ C^{O} = 16, \ L^{N} = 3, \ L^{O} = 7$$

$$u^{N} = \frac{C^{N}}{D'} = \frac{6}{5} > 1$$

NEU & Polyu

A mixed-criticality task

Calculating the mapping



$$C^N = 6, C^O = 16, L^N = 3, L^O = 7, D = 13, D' = 5$$

Using our mapping equation:

$$S_{i}^{N} = \begin{cases} u_{i}^{N} & \text{when } u_{i}^{N} < 1\\ \frac{C_{i}^{N} - L_{i}^{N}}{D_{i}^{\prime} - L_{i}^{N}} & \text{when } u_{i}^{N} \ge 1 \end{cases}$$
$$S_{i}^{O} = \begin{cases} 0 & \text{if task is } LO\\ \frac{C_{i}^{O} - S_{i}^{N}D_{i}^{\prime} - L_{i}^{O}}{D_{i} - D_{i}^{\prime} - L_{i}^{O}} & \text{if task is } HI \end{cases}$$

A mixed-criticality task

Calculating the mapping



$$C^N = 6, C^O = 16, L^N = 3, L^O = 7, D = 13, D' = 5$$

$$S^{N} = \frac{C^{N} - L^{N}}{D' - L^{N}} = \frac{6 - 3}{5 - 3} = \mathbf{1.5}$$
$$S^{O} = \frac{C^{O} - S^{N}D' - L^{O}}{D - D' - L^{O}} = \frac{16 - 1.5 * 5 - 7}{13 - 5 - 7} = \mathbf{1.5}$$



- Mixed-Criticality tasks are encapsulated into mixed-criticality container-tasks
- Scheduling these mixed-criticality container-tasks to physical processors with a partitioned or global algorithm

- Mixed-Criticality tasks are encapsulated into mixed-criticality container-tasks
 - Encapsulating the vertex v0



- Mixed-Criticality tasks are encapsulated into mixed-criticality container-tasks
 - Setting the deadline of the mixed-criticality container-task



- Mixed-Criticality tasks are encapsulated into mixed-criticality container-tasks
 - Encapsulating the vertex v1



- Mixed-Criticality tasks are encapsulated into mixed-criticality container-tasks
 - Encapsulating the vertex v5



- Mixed-Criticality tasks are encapsulated into mixed-criticality container-tasks
 - State transition: normal -> critical - - - - → Additional overload





- Mixed-Criticality tasks are encapsulated into mixed-criticality container-task
 - In critical states





- Mixed-Criticality tasks are encapsulated into mixed-criticality container-tasks
- Scheduling these mixed-criticality container-task to physical processor with a partitioned or global algorithm
 - First, system is in normal state and only normal mixed-criticality container-tasks are scheduled
 - When the system experiences state transition, all normal mixed-criticality container-tasks are banned from scheduling and only critical mixed-criticality container-tasks can run

Outline

Background & Contributions

• Preliminaries

Semi-Federated Mixed-Criticality Algorithm

• Evaluation

- Our task sets are generated based on OpenMP benchmarks and ompTG tool.
- Evaluate the acceptance ratio of task sets with different normalized utilizations.
- We compare our results with those in existing work for federated mixed-criticality scheduling.

Evaluation



Comparison of FMC-Li and SFMC

Semi-federated mixed-criticality algorithm has better schedulability performance

|--|

Thanks for your attention !